Compiling X10 for Scalable High Performance

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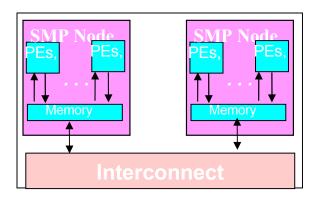
Talk Outline

- What is X10? Why should I care?
- X10 in a Nutshell
- HPC Challenge (Class 2) Results
- Compilation Challenges & Opportunities
- Conclusions

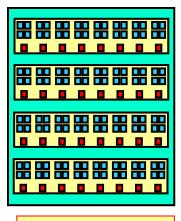
What is X10?

- X10 is a new language developed in the IBM PERCS project as part of the DARPA program on High Productivity Computing Systems (HPCS)
- X10 is an instance of the APGAS programming model in the Java family of languages
- X10 is an open-source project (http://x10-lang.org)

The current architectural landscape



Power5 Clusters

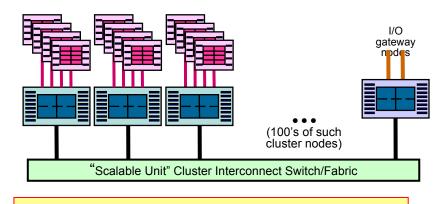


P7 supernode



Blue Gene

Multi-core processors, with accelerators
e.g. Sun Niagara
e.g. Intel multicore, IXP
e.g. IBM Cell
e.g. GPGPUs



Road Runner: Cell-accelerated Opteron

The current architectural landscape

- Substantial architectural innovation is anticipated over the next ten years.
 - Hardware situation remains murky, but programmers need stable interfaces to develop applications
- Heterogenous acceleratorbased systems will exist, raising serious programmability challenges.
 - Programmers must choreograph interactions between heterogenous processors, memory subsystems.

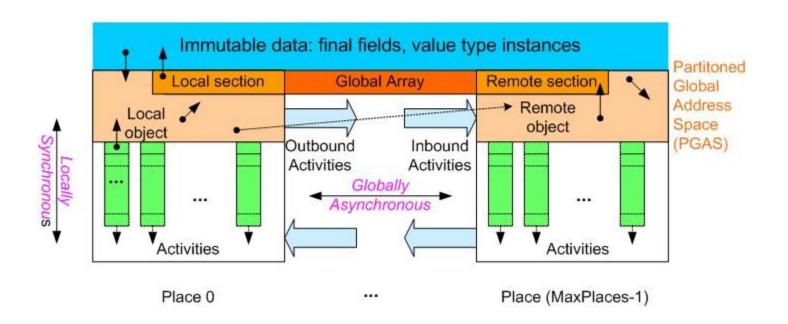
- Multicore systems will dramatically raise the number of cores available to applications.
 - Programmers must understand concurrent structure of their applications.
- ◆Applications seeking to leverage these architectures will need to go beyond dataparallel, globally synchronizing MPI model.
- ◆These changes, while most profound for HPC now, will change the face of commercial computing over time.

Fundamental Challenge

- What is a good Programming Model for these machines?
 - How do we migrate existing users beyond MPI so that they can productively use these machines, specifically for HPC, at scale?
 - How do we make it easy for new classes of users to program such machines?
 - The need for a common programming model has never been more urgent.

X10 Concepts **PGAS:** Replicated Data **Ł**ocal Heap Remote Heap Immutable data: final fields, value type instances Local section Giobal Array Remote section **Locality Rule:** Remote Any access to a mutable Local . object datum must be performed object Synchronous Outbound Inbound by a local activity -> Locally Activities Activities remote data accesses can Globally be performed by creating Asynchronous remote activities Activities Activities Place 0 Place (MaxPlaces-1) **Place** = collection **Activity** = sequential of resident activities computation that runs and objects at a place Storage classes: **Locally Synchronous: Globally Asynchronous:** Activity-local Guaranteed coherence for No ordering of inter-place ◆ Place-local activities → use explicit local heap → sequential Partitioned global consistency synchronization for coherence **Immutable**

X10 Constructs



Fine grained concurrency	Atomicity	Global data-structures
• async S	• atomic S	• points, regions, distributions, arrays
	• when (c) S	distributions, arrays
Place-shifting operations	Ordering	
• at (P) S	• finish S	
	• clock	

Two basic ideas: Places and Asynchrony

Parallel HelloWorld

```
import x10.io.Console;
class HelloWorldPar {
 public static def main(args:Rail[String]):void {
    finish ateach (p in Dist.makeUnique()) {
      Console.OUT.println("Hello World from Place" +p);
(%1) x10c++ -o HelloWorldPar -O HelloWorldPar.x10
(%2) mpirun -n 4 HelloWorldPar
Hello World from Place (0)
Hello World from Place (2)
Hello World from Place (3)
Hello World from Place (1)
(%3)
```

Fibonacci (brute force)

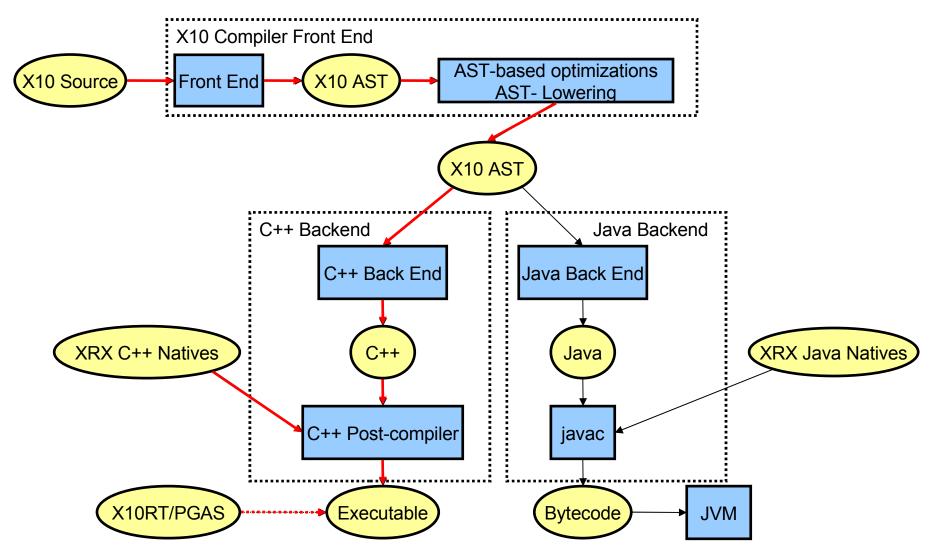
```
public class Fib {
  * Used as an in-out parameter to the computation.
  * When the Fib object is created, r indicates the number to compute.
  * After the computation has completed, r holds the result (Fib(r)).
 var r:int;
 public def run() {
  if (r<2) return; // r already contains Fib(r)
  val f1 = \text{new Fib(r-1)};
  val f2 = new Fib(r-2);
  finish {
   async f1.run();
   f2.run();
  r = f1.r + f2.r;
```

Overview of Features

- Many sequential features of Java inherited unchanged
 - Classes (w/ single inheritance)
 - ☐ Interfaces, (w/ multiple inheritance)
 - ☐Instance and static fields
 - ☐ Constructors, (static) initializers
 - Overloaded, over-rideable methods
 - ☐ Garbage collection
- Structs
- Closures
- Points, Regions, Distributions, Arrays

- Substantial extensions to the type system
 - Dependent types
 - ☐Generic types
 - ☐ Function types
 - ☐ Type definitions, inference
- Concurrency
 - ☐ Fine-grained concurrency:
 - async (p,l) S
 - □ Atomicity
 - atomic (s)
 - □Ordering
 - L: finish S
 - □ Data-dependent synchronization
 - when (c) S

X10 Compilation



X10 Project Status

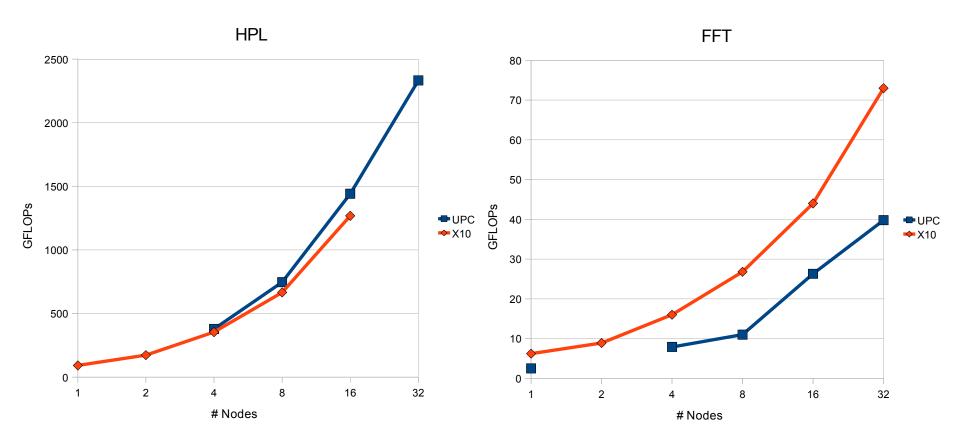
- ◆ X10 is an open source project (Eclipse Public License)
 - Documentation, releases, mailing lists, code, etc. all publicly available via http://x10-lang.org
- XRX: X10 Runtime in X10 (14kloc and growing)
- X10 1.7.x releases throughout 2009 (Java & C++)
- ◆ X10 2.0 will be released this week (rc1 available now)
 - ☐ Java: any platform with Java 5
 Single process (all places in 1 JVM)
 - □C++:

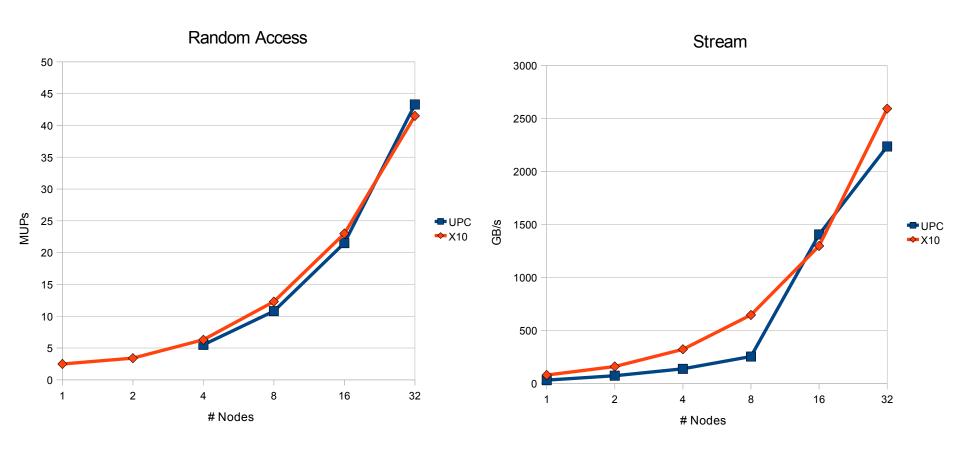
Multi-process (1 place per process)

- aix, linux, cygwin, macos, solaris
- x86, x86_64, PowerPC, Sparc
- x10rt: APGAS runtime (binary only) or MPI (open source)

HPC Challenge Benchmarks

- Data taken from X10/UPC HPCC'09 submission
 - (full details: http://www.x10-lang.org/hpcc09)
- Used Power 5+ Cluster at POK (v20)
 - P575+, 1.9GHz, 16CPUs/node; 64GB DDR2 memory/node; 32 compute nodes, 28 dedicated, 4 shared;gpfs
 - Dual plane HPS switch
 - Rated performance: 7.6GFlops/s per CPU
- In the process of gathering final data for SC'09 BOF





X10 Compilation Challenges

- All of the usual issues with OO languages
 - Virtual/interface dispatch
 - Small methods, class libraries & frameworks
 - _ ...
 - plus closures and higher-order functions
- Concurrency/Communication
 - Recognize idiomatic async/finish patterns reduce async termination traffic
 - Optimize message traffic hoist "loop invariant" messages eliminate unused object fields from messages

Random Access

```
static def runBenchmark(rails: ValRail[Rail[Long]],
  logLocalTableSize: Int, numUpdates: Long) {
  val mask = (1<<logLocalTableSize)-1;</pre>
  val local updates = numUpdates / Place.MAX PLACES;
  finish for ((p) in 0..Place.MAX PLACES-1) {
     async (Place.places(p))
     @Immediate finish {
       var ran:Long = HPCC starts(p*(numUpdates/Place.MAX PLACES));
       for (var i:Long=0; i<local updates; ++i) {
          val place id = ((ran>>logLocalTableSize) & (Place.MAX PLACES-1)) as Int;
          val index = (ran & mask as Int);
          val update = ran;
          val dest = Place.places(place id);
          val rail = rails(place id) as Rail[Long]{self.at(dest)};
          @Immediate async (dest) {
            rail(index) ^= update:
          ran = (ran << 1) ^ (ran<0L ? POLY : 0L);
       } } }
```

X10 Compilation Opportunities

- Exploiting dependent types

 Drive method specialization and loop versioning
- User directed concurrency refactoring, annotationdriven loop transformations, use IDE tooling to enable iterative loop between user & compiler.
- X10 compiled to both C++ and Java
 Neither is always the best choice. Are there interesting things to be learned by studying together?

Conclusions

- X10/APGAS: a programming language/model for multi-core, clusters, accelerators
- Abundance of interesting compilation challenges
- X10 Innovaton Grants
 http://www.ibm.com/developerworks/university/innovation/x10.html
 Short timeline: due 11/25, awarded late 2009/early 2010
 course materials, applications/frameworks/DSLs, tools
- More information on X10: http://x10-lang.org